

A Tour of Computer Systems

Based on content
created by Dr. Steve Ko

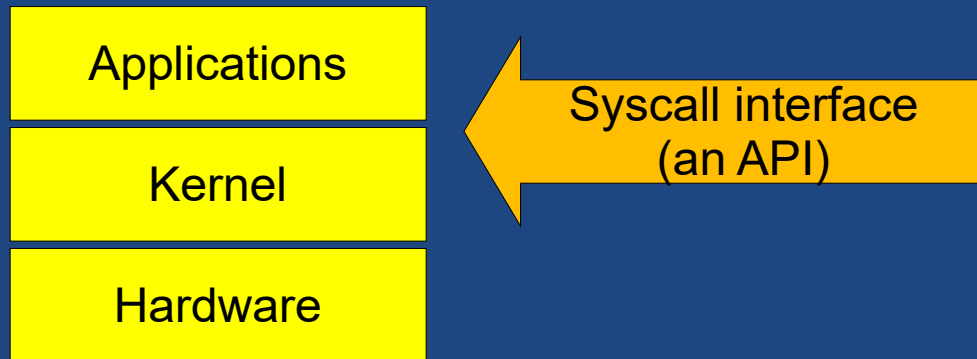
Topics

- 1) For a program to **run**, what is needed?
- 2) How does a computer's **hardware** work?
- 3) What does the **OS Kernel** do?
- 4) How does a **program** interact with the **OS**?

Systems Programming

OS Stack

- Let's discuss the *terminology* necessary for the course and generally for computer systems.
- OS Stack
 - ..



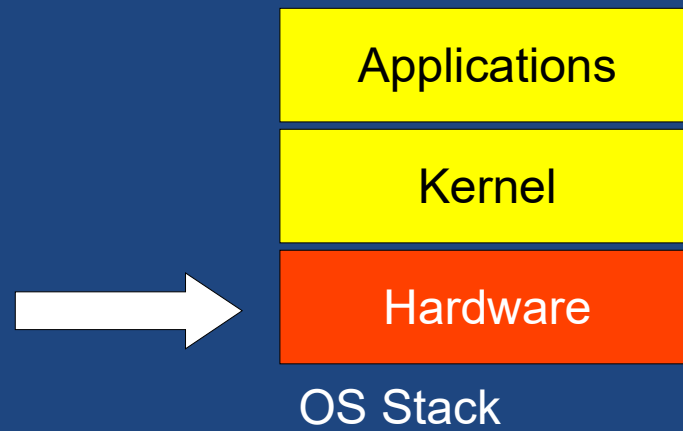
OS Stack

CMPT 201 deals extensively
with the syscall interface

Systems Programming

- Systems programming: ...
 - Low-level languages (e.g., C, C++, Rust) give you the ability to do systems programming, e.g., ..
(Python and Java don't allow you to do that)
- Higher-level programs
 - Don't typically need a systems programming language, unless it needs high performance.
 - Choose a language that fits the target program's goals.
- Let's look at stack bottom up.

Hardware Layer

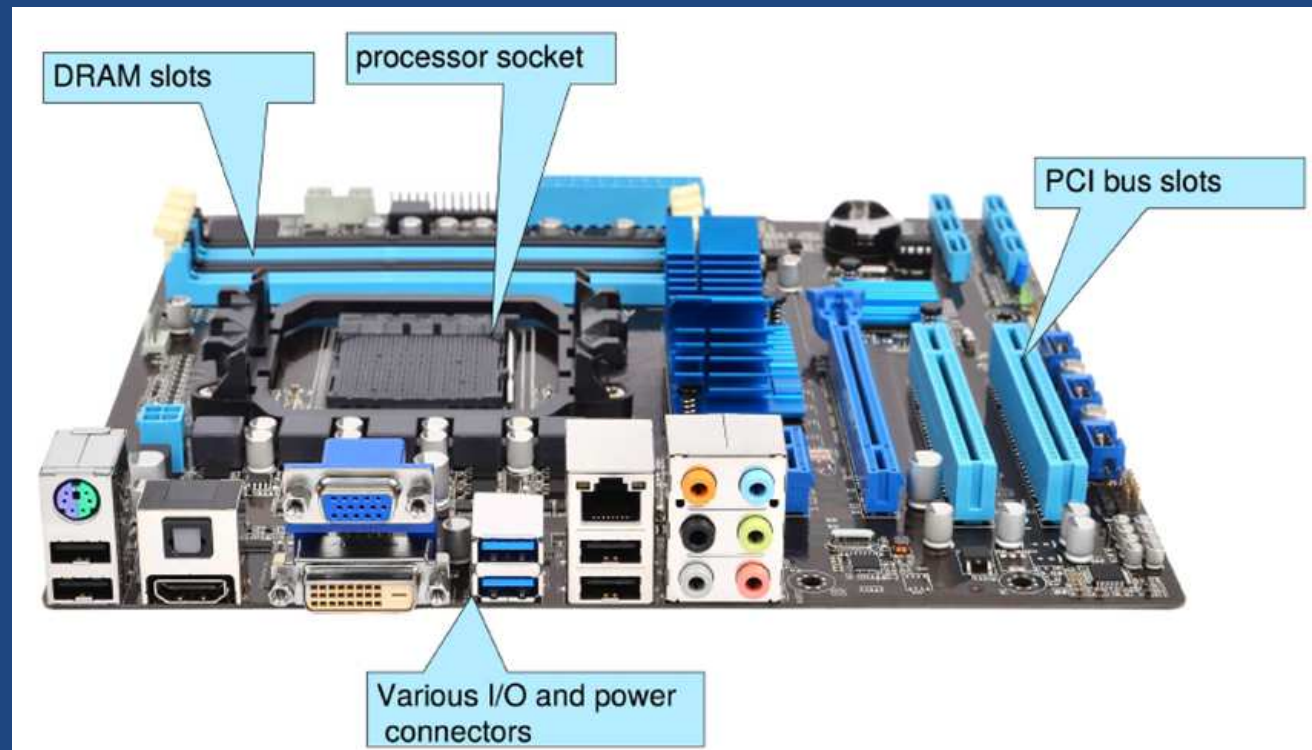


Components in Computing

- 2 Fundamental Components in Computing:
 - ..
Handled by the Central Processing Unit (CPU)
 - ..
Handled by memory (main memory (RAM) and storage)
- E.g., $a + b \Rightarrow c$
 - What is the computation?
 - What is the data?

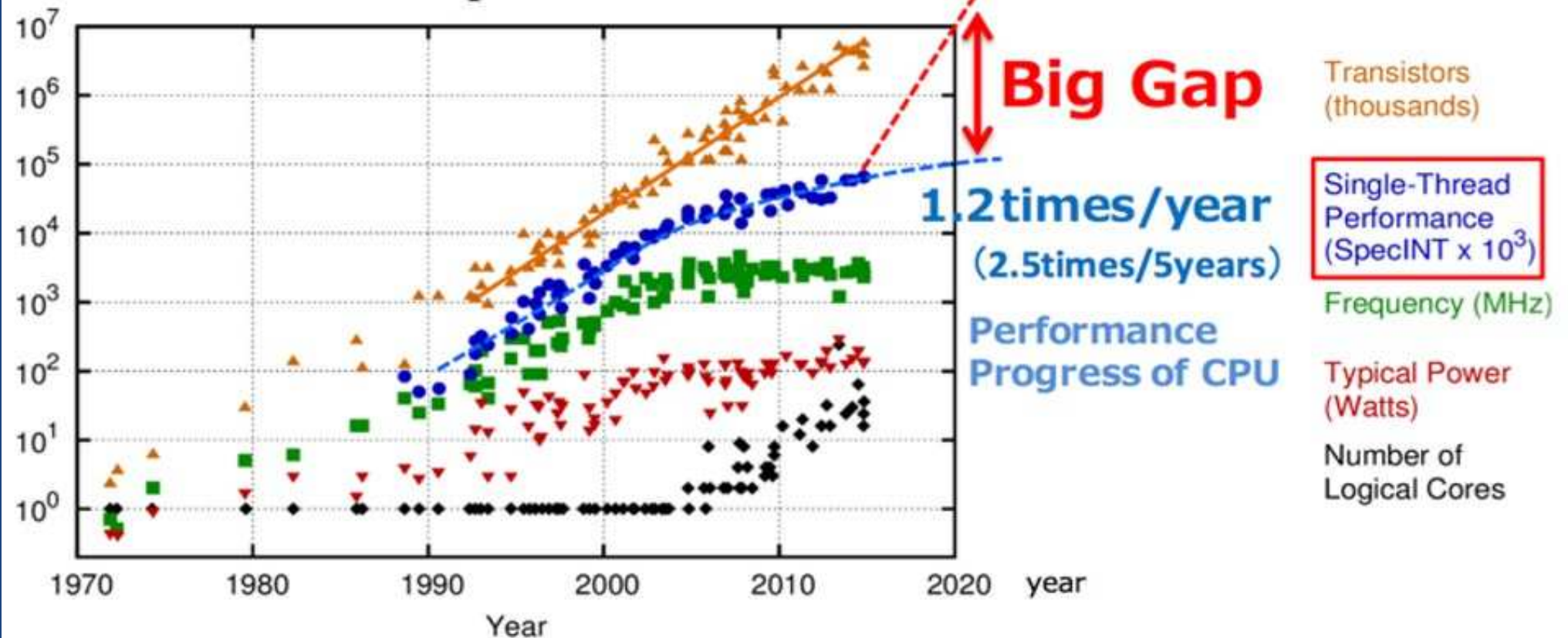
PC Motherboard

- von Neumann architecture
 - Current fundamental model of computer design.
 - Fetch **data** from **memory** to provide to the **CPU**.
- Hardware components:
CPU, **memory**,
and **I/O devices**.



Evolution of CPU: Moore's Law

40 Years of Microprocessor Trend Data



Pre early 2000: frequency $\times 2$ every 18 months
Post 2005: core count $\times 2$ every 18 months

Reference: Ahmet Ceyhan, Interconnects for Future Technology Generations: CMOS with Copper/Low- κ and Beyond, PhD Thesis, 2018

Evolution of Memory

- CPU needs data from memory
 - CPU was getting faster, so memory access had to get faster too.
 - Speed of memory access limited by ..
 - Memory is far away from CPU, and much too slow

CPU

RAM

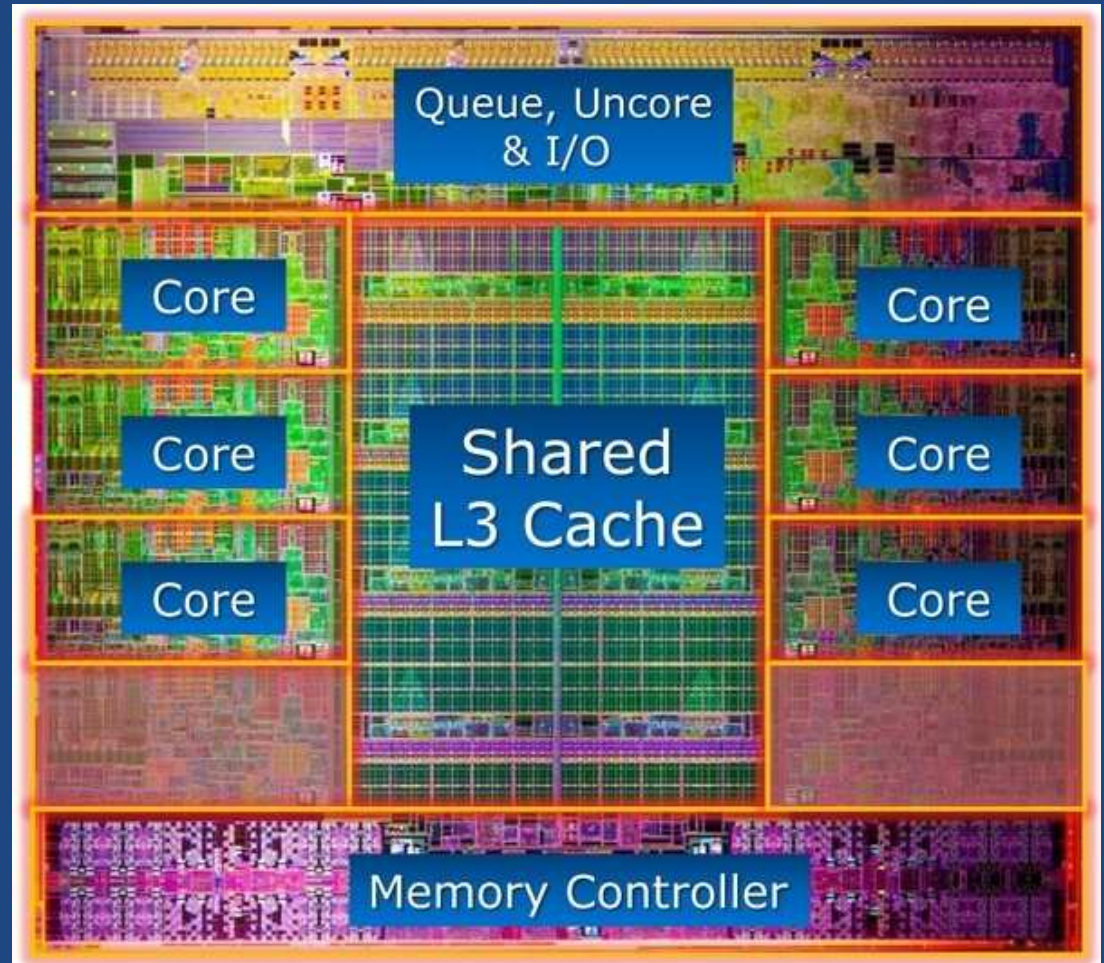


CPU vs Memory Speed

- “Solve” speed gap between CPU and memory access
 - .. very small memory inside a CPU; hold data items from memory. Very close to CPU, so very fast access to data
- Add cache
 - Much larger in size than registers, but much smaller than memory.
 - Quite close (physical distance) to CPU, so..
 - Nowadays processors have many caches:
 - L1 cache ~512 KB (smallest, closest, fastest)
 - L2 cache ~8MB
 - L3 cache ~32MB (large, slowest)

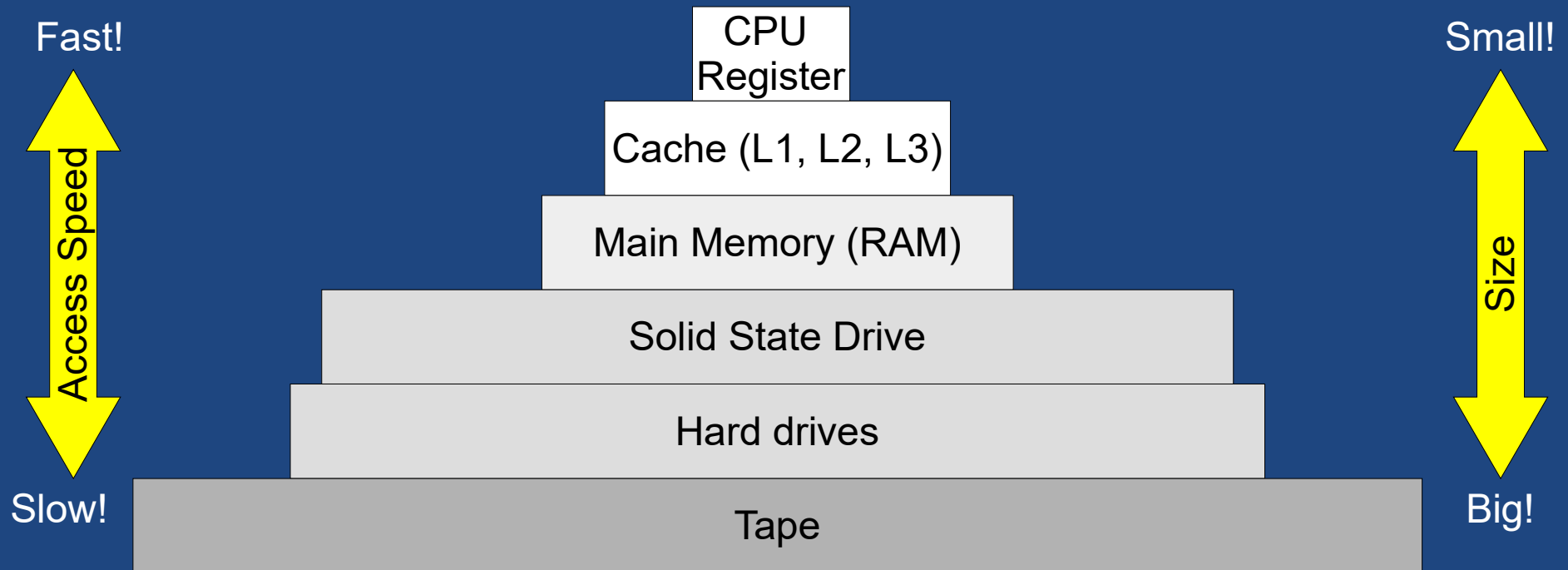
Multi-core Processor

- Desktop CPU today
 - One processor chip
 - Multiple Cores
 - Shared & private caches



Memory Hierarchy

- We want the CPU to feel like it has access to..
 - Intelligently bring data in from large-slow devices (hard drives) into small-fast devices (memory, cache).



Memory Hierarchy

- Trade-offs

- ..
Bigger size typically means more expensive
(size correlates with price).
- ..
faster means closer to CPU.
- ..
"Commit" means moving data from memory to disk;
i.e., changing state of data from temporary to permanent.
 - e.g., `git commit`.
- ..
SSD vs. HDD vs. tape: SSD's fastest but least reliable.
A tape is slowest but most reliable and lasts longer.

CPU Architectures

- Instruction Set Architectures (ISA)
 - ..
 - Compiler translates C programs into machine instructions.
 - E.g. ISAs: x86, ARM, RISK-V ("risk-five")
- 32-bit vs. 64-bit architectures
 - For CMPT 201, we care most about 32-bit vs 64-bit because it..

ABCD - Pointers

- What is a **pointer** in your C program?

- a) A memory address.
- b) A variable storing a memory address.
- c) The data stored in an array.
- d) The address of the current instruction.

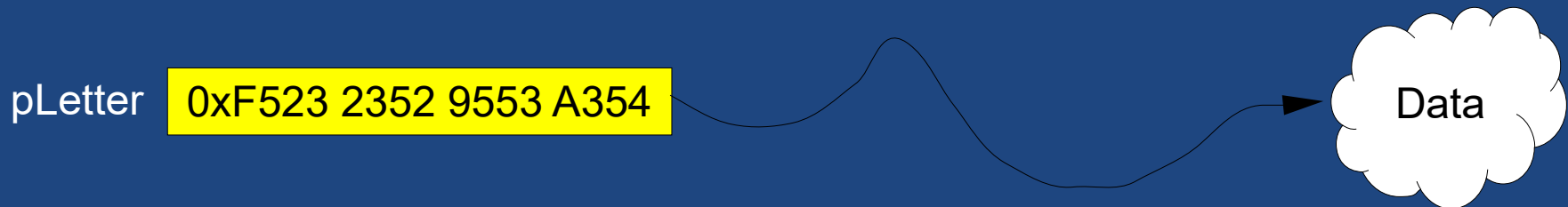
- Which of the following is true about the following code?

```
char* pLetter;  
long long* pCounter;
```

- a) sizeof(pLetter) < sizeof(pCounter)
- b) sizeof(pLetter) > sizeof(pCounter)
- c) sizeof(pLetter) == sizeof(pCounter)
- d) Depends on if the system is 32-bit or 64-bit

32 vs 64 bit Register Size Implications

- **Big Computations:**
In 32-bit, can do 64-bit computation in multiple operations.
- ..
(32-bit uses 32-bit pointers & 64-bit uses 64-bit pointers).



- ..
Pointer size controls the size of the memory address space
- **Bus Width / Memory Channel Width**
Pointer size affects # physical wires connecting to memory.
 - With 64-bits:
need 64 wires to transfer address from CPU to memory.
need 64 wires to transfer data from memory back to CPU

House Analogy

- Imagine the government uses the following form to identify houses in a town:

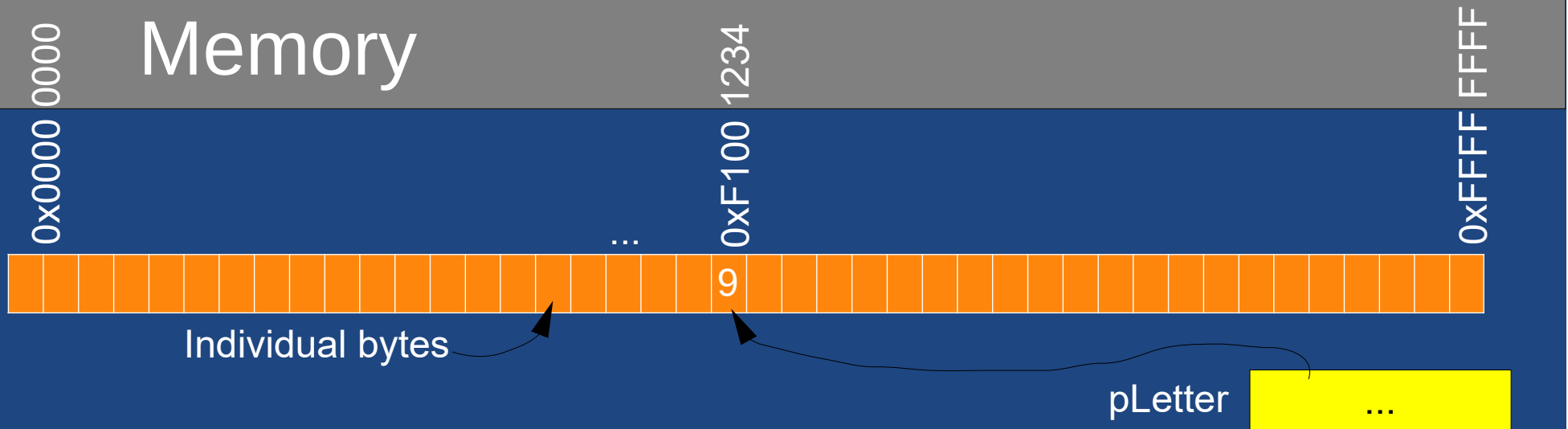
House#: 

Street #: 

- Each house has a house number and a street number (base 10).
- How many total houses can be in the town?

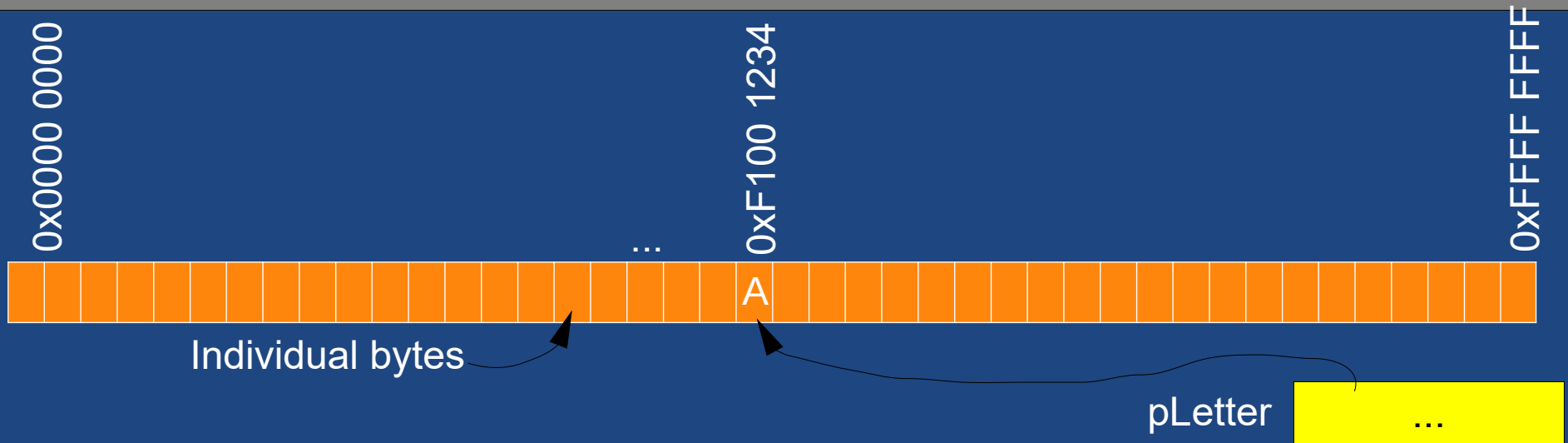
- a) 100,000 Houses
- b) 999,999 Houses
- c) 1,000,000 Houses
- d) 9,999,999 Houses

- If you have 8-bit system, pointers are 8 bits; therefore can address 256 bytes in memory (0-255).



- Memory made up of bytes (1 byte = 8 bits).
 - ..
- 32-bit vs 64-bit Word Size
 - The number of bits stored in a CPU's register.
- In a 32-bit system (32-bit word):
 - Addresses are 32-bits:
0x0000 0000 to 0xFFFF FFFF
 - (Data is retrieved from memory 32-bits at a time (4 bytes) but memory addresses are still byte addresses)

ABCD: Pointer Values



- Which of the following is true?

`char ch = 'A';`

`char* pLetter = &ch`

- a) `pLetter == 'A'`
- b) `pLetter == 0x0000 000A`
- c) `pLetter == 0xF100 1230`
- d) `pLetter == 0xF100 1234`

ABCD - Memory

- Which of the following is true?

a) 1,000 = MB, 1,000,000 = KB, 1,000,000,000 = GB
b) 1,000 = GB, 1,000,000 = MB, 1,000,000,000 = KB
c) 1,000 = KB, 1,000,000 = MB, 1,000,000,000 = GB
d) 1,000 = GB, 1,000,000 = KB, 1,000,000,000 = MB

- If memory (RAM) stored just 16 bytes (16 locations), how many bits do we need in our address?

a) 2-bits
b) 4-bits
c) 8-bits
d) 16-bits

Why 64-bits?

- Why are most computers 64-bit architectures now?

- Registers are 64-bit
- Pointers are 64-bit
- Allows us to..

$$2^{64} = 18,446,744,073,709,551,616$$

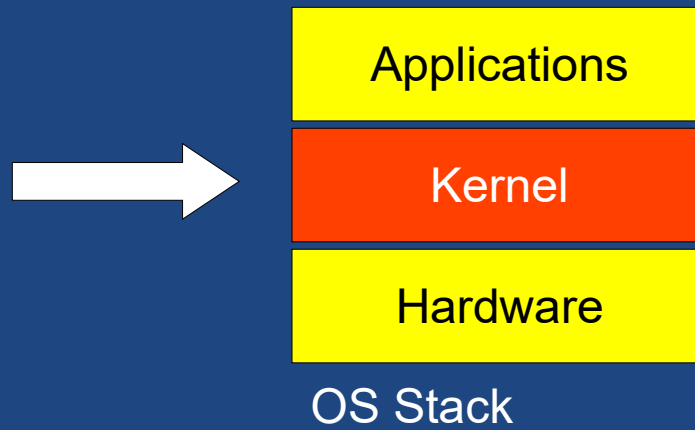
$$\sim 17,000,000,000 \text{ GB} = 16 \text{ Exabytes (VERY large)}$$

$$\begin{aligned} \text{GB} &= 2^{30} \\ &= 1073741824 \end{aligned}$$

- In a 32-bit architecture, how much memory can the CPU access?

- a) 65,526 bytes
- b) 2,147,483,648 bytes
- c) 4,294,967,296 bytes
- d) 18,446,744,073,709,551,616 bytes

Kernel Layer



What is the OS?

- Operating System (OS)

..

- OS Includes:

- ..

- Main part that actively manages resources.

- Supporting tools:

- such as GUI, command line;

- These are what differentiates Linux distributions (“distros”)

What does a Kernel do?

- Kernel's Role

- ..

- many programs want to access the hardware at the same time
 - kernel manages (mediates) access

- ..

- the kernel controls programs (running, stopping, etc.).

- ..

- the kernel provides protection (isolation) for users and programs.

- E.g., users can't access each other's data
 - E.g., programs can't interfere with each other's execution.

Event-Driven

- When does a kernel do some work?
 - Generally, the OS lets other programs run and waits for something it needs to do.
 - The kernel is..
It responds to events.
- Events can be:
 - ..
a hardware event like a mouse click, or network packet received
 - ..
a user-space-application generated call to the kernel
e.g., application asking kernel to `printf()` to the screen.
 - ..
a software interrupt that announces an event to a process
e.g., `SIGINT` = ctrl+c, `SIGSEGV` = segmentation (page) fault

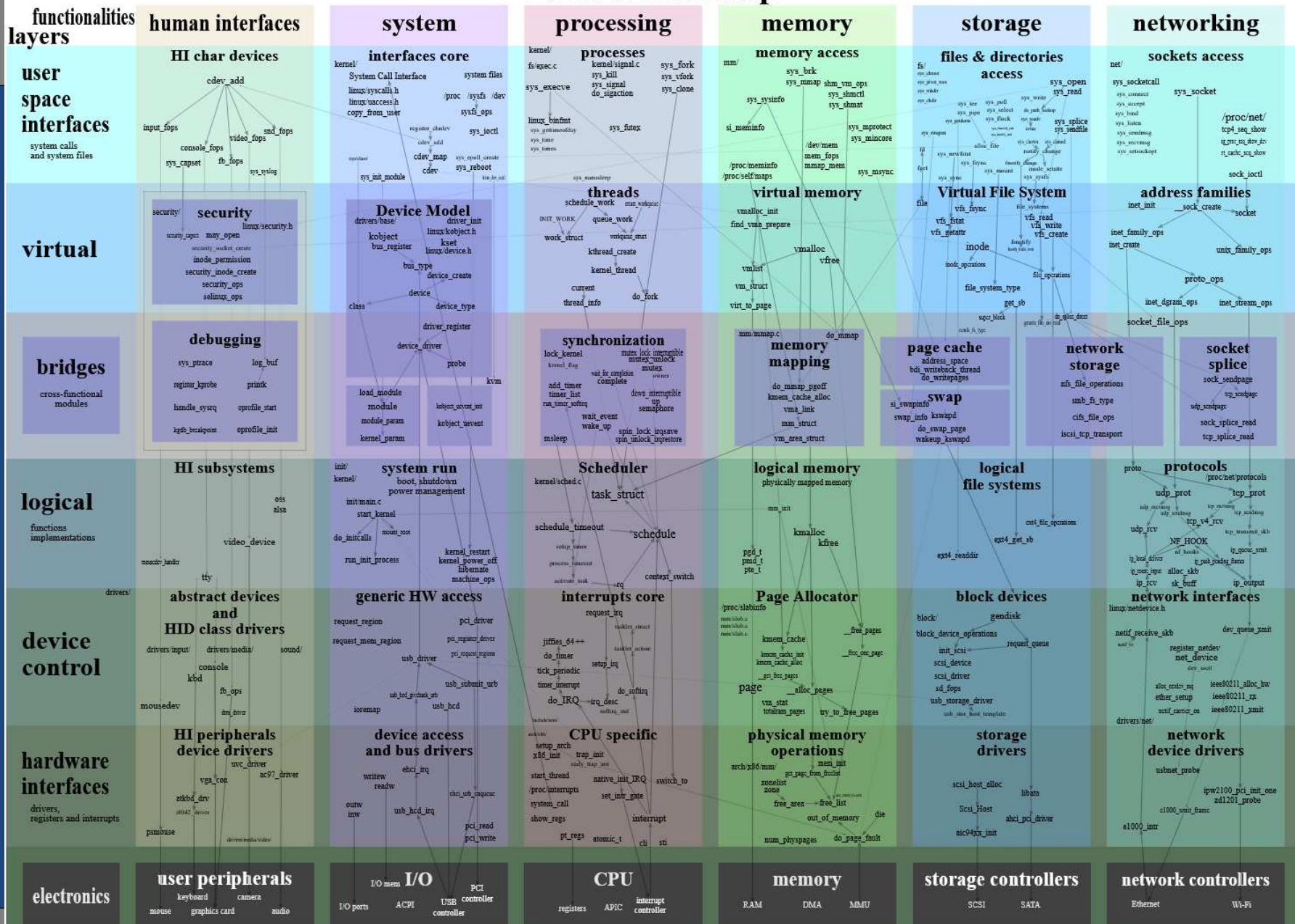
User Mode vs. Kernel Mode

- **Privilege mode of CPU execution**
 - Kernel Mode runs the OS kernel; allows full privilege and full access to the hardware. Often called "Ring 0"
 - User Mode runs applications; ..
 - E.g., instructions that allow direct access to hardware
 - E.g., access to certain regions of memory (kernel memory)
- Modern CPUs run in one of those two modes at a given moment.
- **ABCD**: Which best explains why we need a user mode?
 - a) Isolation
 - b) Efficiency
 - c) Null pointers
 - d) Abstraction

Root user (aside)

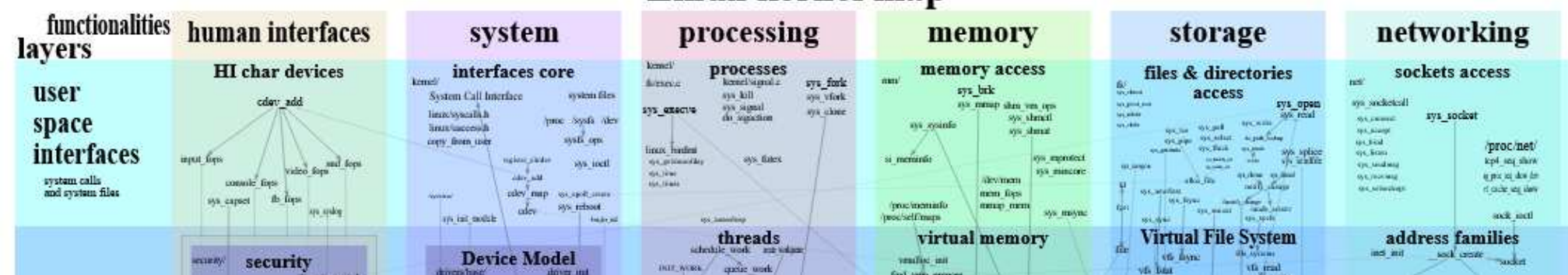
- User / Kernel Mode vs Root User
 - The “mode” (privilege level of code) is different than the user-level
 - The root user is still a user, but with full admin privileges
 - Root can run programs and access files that normal users cannot.
 - Root user often called a super user.
 - Root user cannot access kernel memory or protected instructions.

Linux kernel map



Important Terms in the Kernel

Linux kernel map



- **System**

- .. : every device needs a device driver to control it.
E.g., network card device driver talks to hardware to send/receive data to/from the physical network.

- **Processing**

- Processes, threads, synchronization, and scheduling

Covered later

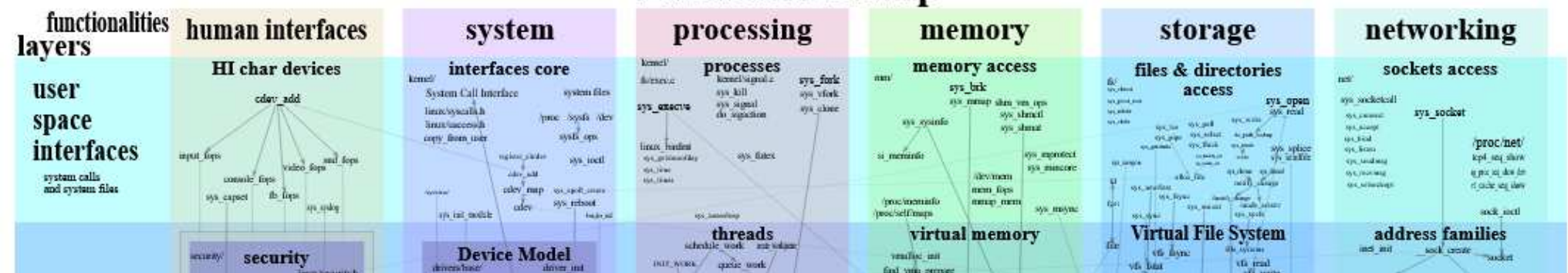
- **Memory**

- Virtual memory, physical memory, and paging

Covered later

Important Terms in the Kernel (cont)

Linux kernel map



- **Storage**

- File systems, and VFS (Virtual File System).

VFS is an interface:

..

e.g., read and write.

- By looking like a normal file, many tools can seamlessly work with it
e.g., `cat /proc/cpuinfo`

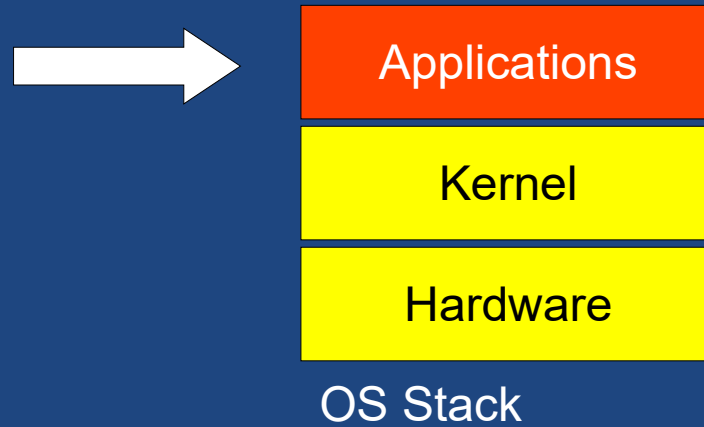
- **Networking**

- Sockets, TCP, UDP, and IP

Covered later

- Which of the following is true?
 - a) The root user runs programs in kernel mode.
 - b) Syscalls allow the kernel to execute user-level applications.
 - c) A hardware interrupt is generated when dereferencing a null pointer.
 - d) User mode prevents applications from executing privileged instructions.

Applications Layer



Lifetime of a Program

(briefly)



Source Code



Compilation

Executable
(machine code)



Memory Loading

Running Program

Compilation vs. Interpretation

(briefly)

- Two major ways to run a program:
 - Compilation (e.g., C, C++)
 - Interpretation (e.g., Python, Bash)
- Performance vs Portability Trade-off
 - Compilation has better performance:
it directly generates machine code to execute.
 - ..
machine code for one specific ISA
E.g., can't run AMD64 executable on ARM64 machine
 - Interpretation is slower, but same script can run anywhere
there is an interpreter.

POSIX

(briefly)

- **POSIX = ..**
 - A standard for (user-level) software portability across different OSs.
 - Includes programming interface (file I/O, C standard library, etc.) and shell utilities
 - We see it in C to: specifies what features we need:
`#define _POSIX_C_SOURCE 200809L`



ABI

- ABI = ..
- Similar to API = ..
 - An API is at the code level:
Your code calls or accesses the functions of the API, such as provided by a library.
 - An ABI is an interface for a binary (an executable) that an OS defines.
- Compilers generate executables that follow the ABI for the OS
 - E.g., Windows ABI is different from Linux ABI.
Cannot copy a Windows binary (`.exe`) to a Linux machine and run it (and vice versa).

Virtualization

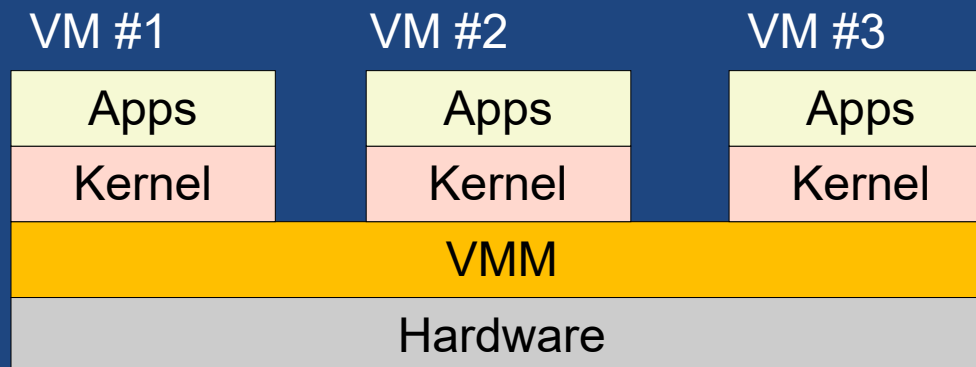


Virtualization of Traditional OS Stack

- Virtualization allows..
 - Lets us be much more flexible!
 - Software can control the environment:
"Spin up 3 virtual machines to host new databases"
- ..
software that *provides* virtualization.
 - Also called the Virtual Machine Monitor (VMM)
 - Hypervisor can run at different levels of our OS stack, giving different levels of flexibility

On Hardware

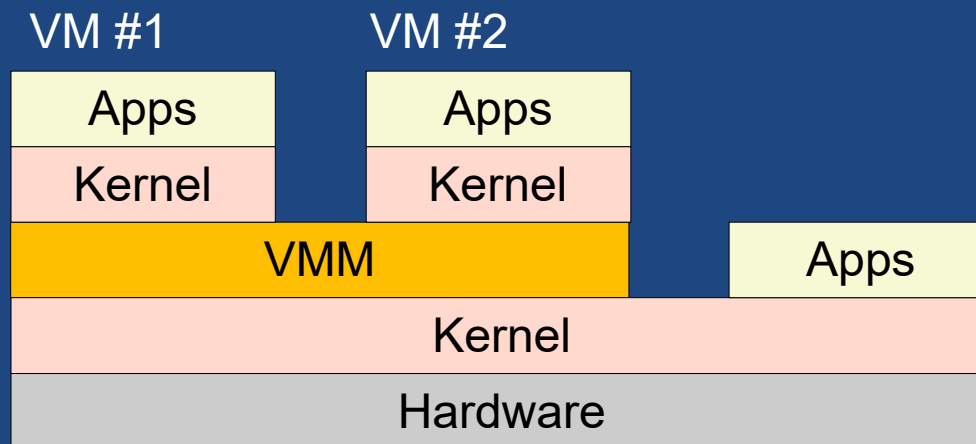
- VMM Directly atop Hardware
 - VMM..
 - This is often used in a data center environment.



On Kernel

- VMM atop the Kernel

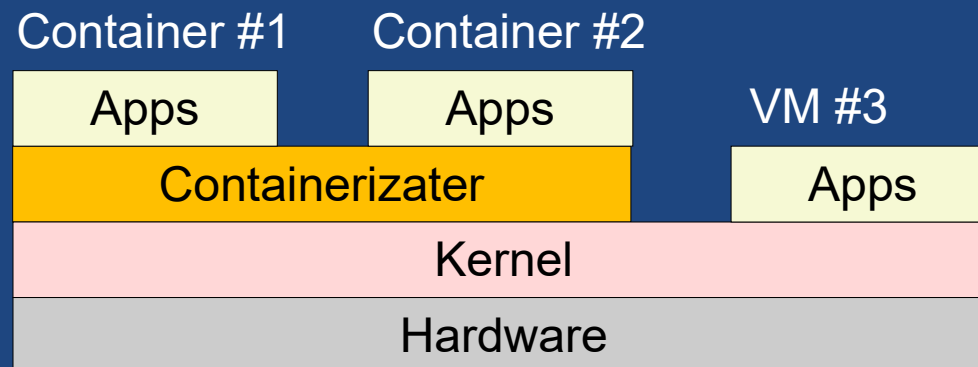
- A VMM is an application running atop a kernel, along with other applications.
- The VMM creates/runs/manages VMs.
- This is often used in a desktop environment, e.g., VMWare Workstation, VirtualBox, QEMU.



Containerization

- **Containerization**

- Containerization creates a **container** not a virtual machine.
- Container includes..
- Uses the same OS kernel as rest of the system
- Uses Linux features for isolation: **process isolation** (namespaces), **resource control/isolation** (cgroups), etc.
- This is the most popular form of virtualization these days, e.g., Docker, Podman.



ABCD - Virtualization

- Which of the following is a major benefit of **virtualization**?
 - a) Allows user level applications to call the kernel.
 - b) Allows parts of the OS stack to be swapped out under software control.
 - c) Allows the kernel to control different pieces of hardware when they are connected at runtime.
 - d) Allows application to run without using an OS kernel.

Summary

- OS Stack is the layers of service
 - Hardware, Kernel, Application.
- Memory hierarchy
 - allows programs to access large memories quickly
- Pointers hold addresses,
 - 32 vs 64 bits limit how much memory we can access
- Kernel mode gives OS kernel access to all resources
 - User mode limits what an application can do.
- Applications use the OS's ABI to use services
- Virtualization allows parts of the OS stack to be swapped out under software control.